What is a Database?

- A **database** is any collection of data.

- A **DBMS** is a software system designed to maintain a database.

- We use a DBMS when
  - there is a large amount of data
  - security and integrity of the data are important
  - many users access the data concurrently
Consider a Phone Company, such as AT&T

Kinds of information they deal with:

- customer records
- employee records
- billing information
- management records
- switching and wiring diagrams
- customer service orders
With all that data, AT&T must be concerned with questions such as:

- Where is the information kept?
- How is the data structured?
- How is the data kept consistent?
- How is the data described?
- How is the data kept secure?
- How do different pieces of data interrelate?
Without a DBMS, we'd have:

- Access by a collection of ad hoc programs in C++, Java, PHP, etc.
- Data stored as bits on disks organized as files
- Users of the data

There is no control or coordination of what these programs do with the data
With a DBMS, we have:

- DBMS applications
- Data stored as bits on disks organized as files
- Users of the data

DBMS provides control and coordination to protect the data.
DBMS Languages

DML: data manipulation language
QL: query language
GPL: general purpose languages

DDL: data definition language

query processor
security manager
concurrency manager
index manager

application program(s)

users of the data

system configuration languages
Advantages of Using a DBMS

• Anything you can do with a DBMS, you can do with a file system, a network and a heap of C code

• So why spend the money to buy a DBMS?
  - there is a well defined collection of capabilities common to a certain class of applications
  - for applications in this class, the DBMS already has these capabilities and probably does them better than you could with home-brewed code
Database
Functionality
Advantages of Using a DBMS

Users of the data application program(s) can be provided with many capabilities in a generic way.
Persistence

- A DBMS provides persistent objects, types and data structures
  - *persistent* = having a lifetime longer than the programs that use the data
  - any information that fits the data model of a particular DBMS can be made persistent with little effort
  - *data model* = concepts that can be used to describe the data
Concurrency

- A DBMS supports access by concurrent users
  - *concurrent* = happening at the same time
  - concurrent access, particularly writes (data changes), can result in inconsistent states (even when the individual operations are correct)
  - the DBMS can check the actual operations of concurrent users, to prevent activity that will lead to inconsistent states
Access Control

- A DBMS can restrict access to authorized users
  - security policies often require control that is more fine-grained than that provided by a file system
  - since the DBMS understands the data structure, it can enforce fairly sophisticated and detailed security policies
    - on subsets of the data
    - on subsets of the available operations
Redundancy Control

- A DBMS can assist in controlling redundancy
  - redundancy = multiple copies of the same data
  - with file storage, it's often convenient to store multiple copies of the same data, so that it's "local" to other data and applications
  - this can cause many problems:
    - wasted disk space
    - inconsistencies
    - need to enter the data multiple times
Complex Semantics

- A DBMS supports representation of complex relationships and integrity constraints
  - the semantics (meaning) of an application often includes many relationships and rules about the relative values of subsets of the data
  - these further restrict the possible instances of the database
  - relationships and constraints can be defined as part of the schema
Backup and Recovery

- A DBMS can provide backup and recovery
  - **backup** = snapshots of the data particular times
  - **recovery** = restoring the data to a **consistent** state after a system crash
  - the higher level semantics (relationships and constraints) can make it difficult to restore a consistent state
  - **transaction analysis** can allow a DBMS to reconstruct a consistent state from a number of backups
Views and Interfaces

• A DBMS can support multiple user interfaces and user views
  
  • since the DBMS provides a well-defined data model and a persistent data dictionary, many different interfaces can be developed to access the same data

  • data independence ensures that these UIs will not be made invalid by most changes to the data

  • new user views can be supported as new schemas defined against the conceptual schema
Advantages of Using a DBMS

- persistent objects, types and data structures
- control of concurrent users
- controlling of redundancy
- restricting access (security)
- representation of complex relationships and integrity constraints
- backup and recovery
- multiple user interfaces and user views
Database Users and Roles
DBMS Roles

- Application developers
- Users of the data
- DBMS system developers
- Database designer
- Data definition processor
- Data dictionary
- Data
- Query processor
- Security manager
- Concurrency manager
- Index manager
- System administrator (and DB designer)
DBMS Roles

• Actors On the Scene
  (people interested in the actual data):
  • database administrators
  • database designers
  • systems analysts and application programmers
  • end users
Database Administrators

- acquiring a DBMS
- managing the system
- acquiring HW and SW to support the DBMS
- authorizing access (security policies)
- managing staff, including DB designers
Actors on the Scene

- Database Designers
  - identifying the information of interested in the Universe of Discourse (UoD)
  - designing the database conceptual schema
  - designing views for particular users
  - designing the physical data layout and logical schema
  - adjusting data parameters for performance
Actors on the Scene

- Systems Analysts and Application Programmers (generic database developers)
  - provide specialized knowledge to optimize database usage
  - provide generic (canned) application programs
Actors on the Scene

End Users

- **casual users**: ad-hoc queries
- **naïve or parametric users**: canned queries such as menus for a phone company customer service agent
- **sophisticated users**: people who understand the system and the data and use it in many novel ways
- **standalone users**: people who use personal easy-to-use databases for personal data
DBMS Roles

- **Actors Behind the Scene:**
  people who maintain the environment but aren't interested in the actual data

- DBMS designers and implementers
- tools developers
- operators and maintenance personnel
- database researchers
Actors Behind the Scene

- DBMS designers and implementers
  - work for the company that supplies the DBMS (i.e. Microsoft, Oracle, Sybase, MySQL ...)
  - programmers and engineers
  - design and implement the DBMS
Actors Behind the Scene

- Tools Developers
  - design and implement DBMS add-ons or plug-ins
  - may work for DBMS supplier or be independent
  - kinds of tools: database design aids, performance monitoring tools, user and designer interfaces
Operators and maintenance personnel

- run and maintain the computer environment in which a DBMS operates
- probably work for the database administrator (DBA)
Actors Behind the Scene

- Database Researchers
  - academic or industrial researchers
  - develop new theory, new designs, new data models and new algorithms to improve future database management systems
Data Organization
Schemas and Instances

- A *database instance* is the collective values of all database objects at some point in time
  - also called the *(data) instance* or *(database) state*

- A *schema* describes the database and defines the possible instances
  - also called the *data definition, data dictionary, or meta-data*
Conceptual Data Models

• A *data model* describes the possible schemas (essentially the *meta-schema*)

• A DBMS is designed around a particular data model
  - this is what allows all system components (and humans) to understand the schema and data

• possible data models
  - *relational*,
    - object-oriented, object-relational, entity-relationship,
      - semantic, network, hierarchical, etc.
Physical Data Models

- A *physical data model* describes the way in which data is stored in the computer
  - typically only of interest to database designers, implementers and maintainers ... not end users
  - must provide a well-defined structure that can be mapped to the conceptual schema
  - allows optimization strategies to be defined generically
Three-Schema Architecture

- External View
- External View
- External View

Conceptual Schema

Internal Schema

- generic view
- physical view
- user-specific views
Data Independence

- **physical data independence**
  - conceptual and external schema are defined in terms of the data model, rather than the actual data layout
  - ensures that conceptual and external schemas are not affected by changes to the physical data layout

- **logical data independence**
  - ensures that changes to the conceptual schema don't affect the external views
  - (this is not always achievable)
Transactions
**Transactions**

- *transaction* = an indivisible unit of data processing

- All transactions must have the **ACID** properties:
  - **Atomicity**: all or nothing
  - **Consistency**: no constraint violations
  - **Isolation**: no interference from other concurrent transactions
  - **Durability**: committed changes must not be lost due to any kind of failure
Atomic Transactions

- Fred wants to move $200 from his savings account to his checking account.

1) Money must be subtracted from savings account.
2) Money must be added to checking account.

If both happen, Fred and the bank are both happy.
If neither happens, Fred and the bank are both happy.

If only one happens, either Fred or the bank will be unhappy.

Fred’s transfer must be all or nothing.
Transactions are Atomic

- Transactions must be *atomic* (indivisible)
  - the DBMS must *ensure* atomicity
  - everything happens, or nothing happens
  - boundaries of transaction (in time) are generally set by the application ...

  the DBMS has no means of determining the intention of a transaction
Correct Transactions

- Wilma tries to withdraw $1000 from account 387.

<table>
<thead>
<tr>
<th>No</th>
<th>PIN</th>
<th>Balance</th>
</tr>
</thead>
<tbody>
<tr>
<td>101</td>
<td>8965</td>
<td>10965.78</td>
</tr>
<tr>
<td>387</td>
<td>6643</td>
<td>652.55</td>
</tr>
<tr>
<td>543</td>
<td>4287</td>
<td>8720.12</td>
</tr>
</tbody>
</table>

**constraint:**
account.Balance must be non-negative
any transaction withdrawing more than $652.55 from acct 387 will violate this constraint

Wilma’s transaction cannot be accepted.
Transactions are **Consistent**

- A transaction must leave the database in a **valid or consistent state**
  - valid state == no constraint violations
- A *constraint* is a declared rule defining specifications of database states
- Constraints may be violated temporarily ... but must be corrected before the transaction completes
Fred is withdrawing $500 from account 543.
Wilma’s employer is depositing $1983.23 to account 543.
These transactions are happening at the same time.

Combined result of both transactions must be correct
Transactions are Isolated

- If two transactions occur at the same time, the cumulative effect must be the same as if they had been done in isolation
  
  \[(8720.12 - 500) + 1983.23 = 10233.35\]
  
  \[(8720.12 + 1983.23) - 500 = 10233.35\]

- Ensuring isolation is the task of concurrency control
Durable Transactions

- Wilma deposits $50,000 to account 387.
- Later, the bank’s computer crashes due to a lightning storm.

Wilma’s deposit cannot be lost.
Transactions are Durable

- Once a transaction's effect on the database state has been committed, it must be permanent.

- The DBMS must ensure persistence, even in the event of system failures.

- Sources of failure:
  - computer or operating system crash
  - disk failure
  - fire, theft, power outage, earthquake, operator errors, ...
Transactions

- transaction = an indivisible unit of data processing

- All transactions must have the **ACID** properties:
  - **Atomicity**: all or nothing
  - **Consistency**: no constraint violations
  - **Isolation**: no interference from other concurrent transactions
  - **Durability**: committed changes must not be lost due to any kind of failure